

MATCHING ADJACENT CARDS

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ABSTRACT. In a well-shuffled deck of cards, what is the probability that somewhere in the deck there are adjacent cards of the same rank? What is the average number of adjacent matches? What is the probability distribution for the number of matches? We answer these and related questions for both the standard 52-card deck with four suits and 13 ranks and for generalized decks with k suits and n ranks. We also determine the limiting distribution as n goes to infinity with k fixed.

Here is my offer. First pay a dollar to play the game. Then take a shuffled deck of cards and fan out the entire deck face up onto the table. For each occurrence of adjacent cards with the same value I will pay you a dollar. If there is one match, then you break even. If there are no matches, then you lose your dollar. If there are two or more matches, then you come out ahead. Should you accept my offer?

Get a deck of cards and give it a try. Play at least five times, making sure to shuffle the deck several times before each round.

Playing a few rounds should convince you to take the bet. But would you be willing to pay two dollars to play? What is the maximum you should consider paying? And what is the probability of losing your dollar? We get a better idea with some computer simulation of a hundred thousand games.

In the appendix, you can find code in Mathematica and in Sage to do this. Note that three cards of the same rank occurring consecutively count as two matches, and all four cards of one rank occurring consecutively count as three matches. The maximal number of matches possible is 39, the probability of which is exceedingly small—approximately $6.76672 \cdot 10^{-41}$. (Exercise: find the exact answer.)

For one particular run with 100,000 shuffles, the match count histogram is shown in Figure 1. Although you cannot see this in the histogram, the maximal count that occurred is 13, and that happened just once. The count of 11 occurred 12 times and the count of 12 occurred twice. For this run the average was 3.00125, and so you might be willing to pay up to about three dollars per game. Also, the probability of no matches appears to be between 4% and 5%. The most frequent match counts are 2 and 3, each occurring more than 20% of the time. It is fairly easy to find the expected number of matches using indicator random variables. Let M_i be 1 if card i and card $i + 1$ match and 0 otherwise. Then the number of matches M is the sum $M_1 + M_2 + \cdots + M_{51}$. Expectation is linear, so $E(M) = E(M_1) + E(M_2) + \cdots + E(M_{51})$, but $E(M_i)$ is the same for all i . Finally, $E(M_i) = 3/51$

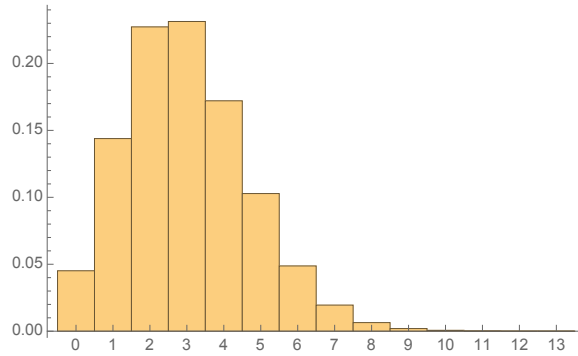


FIGURE 1. A histogram displaying the results of playing our card-matching game 100,000 times.

because there are 51 possibilities for card $i + 1$, and three of them match card i . Therefore, $E(M) = 3$. This easily generalizes:

Theorem 1. *In a randomly arranged deck of kn cards consisting of k suits with n cards in each suit, the expected number of matches is $k - 1$.*

Proof. For $1 \leq i \leq kn - 1$, let M_i be the indicator random variable for the event that cards i and $i + 1$ match, and let $M = \sum M_i$. Then $E(M_i) = (k - 1)/(kn - 1)$ because there are $k - 1$ cards of the same rank as card i and $kn - 1$ possibilities for the next card. Therefore,

$$E(M) = \sum_{i=1}^{kn-1} E(M_i) = \sum_{i=1}^{kn-1} \frac{k-1}{kn-1} = k-1.$$

□

Having found the expected number of matches, we move on to the general question of the complete probability distribution for M . With a standard 52-card deck the number of matches is at most 39 since each rank contributes at most three matches. For general k and n , the number of matches is at most $kn - n$.

Question. For a well-shuffled deck consisting of k suits and n ranks, what is the probability that there are exactly r matches for $r = 0, 1, 2, \dots, kn - n$?

We need some basic definitions. First, we define a card deck with k suits and n ranks to be the multiset of size kn made up of k copies of each integer from 1 to n . For example, for $k = 3$ and $n = 4$ the deck is $\{1, 1, 1, 2, 2, 2, 3, 3, 3, 4, 4, 4\}$. Now we no longer see different suits; the aces are all identical. The number of distinct permutations is no longer $(kn)!$ because the cards in each of the n ranks can be permuted among themselves in $k!$ ways without changing what we see. Therefore, the number of permutations is $(kn)!/(k!)^n$. Letting α_r be the number of permutations with exactly r matches we have

$$P(M = r) = \frac{\alpha_r}{(kn)!/(k!)^n} = \frac{(k!)^n}{(kn)!} \alpha_r.$$

We now focus on the problem of counting the number of permutations with r matches for general k and n . We will deal with $k = 4$ in detail before tackling arbitrary values of k .

Counting permutations with no matches (i.e., $r = 0$) has been of particular interest. Eriksson and Martin [2] named them “Carlitz permutations” and have enumerated them for $k = 2, 3, 4$. The case with $k = 2$ is similar to a linear variation of the relaxed version of the “Problème des ménages.” The classic version asks for the number of ways that n couples can be seated at a circular table so that the men and women alternate and no one sits next to his or her partner. The relaxed version [1] drops the requirement that the genders alternate. Now suppose that they are to be arranged in a single line rather than in a circle. Then the number of possible linear arrangements is 2^n times α_0 , the number of Carlitz permutations with two suits. (See sequences A114938 and A007060 in the On-Line Encyclopedia of Integer Sequences (OEIS) at <https://oeis.org>.)

1. FOUR SUITS ($k = 4$)

The deck is the multiset of size $4n$ consisting of four copies of each number 1 through n , and the number of permutations is $(4n)!/(4!)^n$. Let α_r be the number of permutations with exactly r matches and let

$$A(x) = \sum_{r=0}^{3n} \alpha_r x^r$$

be the associated generating function. In order to use the principle of inclusion-exclusion we define another generating function $B(x) = \sum \beta_m x^m$, which is related to $A(x)$ by the functional equation $A(x) = B(x - 1)$ and with the virtue that the β_m are easier to calculate than the α_r . Essentially, β_m counts permutations that have at least m matches.

One rank with four cards x, x, x, x will account for up to three matches. We can be sure that this rank contributes at least one match by glueing two of the cards together: xx, x, x . To force at least two matches, we can either glue two pairs together as xx, xx or glue three cards together as xxx, x . To force three matches we glue all four cards together: $xxxx$. Now we choose a pattern for each of the n ranks. Let s, t, u, v, w denote the number of ranks having each of the five possible patterns, according to this scheme:

$$\begin{array}{ll} s & x \cdot x \cdot x \cdot x \\ t & xx \cdot x \cdot x \\ u & xxx \cdot x \\ v & xx \cdot xx \\ w & xxxx \end{array}$$

For a given choice of s, t, u, v, w (such that $s + t + u + v + w = n$) the number of objects being permuted is no longer $4n$, but rather $4s + 3t + 2u + 2v + w$. The pattern $x \cdot x \cdot x \cdot x$ has four identical objects, and the patterns $xx \cdot x \cdot x$ and $xx \cdot xx$ each have two identical objects,

and so we divide by the redundancy factor $(4!)^s(2!)^t(2!)^v$. This construction gives a list of

$$(1) \quad \frac{(4s + 3t + 2u + 2v + w)!}{(4!)^s(2!)^t(2!)^v}$$

permutations. A permutation with these parameters has at least $t + 2u + 2v + 3w$ matches. Define β_m for $m = 0, 1, \dots, 3n$ to be the sum of the cardinalities of all the lists such that $m = t + 2u + 2v + 3w$. Thus,

$$(2) \quad \begin{aligned} \beta_m &= \sum_{\substack{s+t+u+v+w=n \\ t+2u+2v+3w=m}} \binom{n}{s, t, u, v, w} \frac{(4s + 3t + 2u + 2v + w)!}{(4!)^s(2!)^t(2!)^v} \\ &= \sum_{\substack{s+t+u+v+w=n \\ t+2u+2v+3w=m}} \binom{n}{s, t, u, v, w} \frac{(4n - m)!}{(4!)^s(2!)^t(2!)^v}. \end{aligned}$$

The permutations with exactly m matches occur once in all of the permutations counted by β_m , but the permutations with more than m matches occur many times. For $r > m$, a permutation with r matches is generated $\binom{r}{m}$ times because each choice of m of its r matches corresponds to a different set of parameters s, t, u, v, w . Therefore,

$$\beta_m = \sum_{r \geq m} \binom{r}{m} \alpha_r,$$

and

$$\begin{aligned} B(x) &= \sum_{m=0}^{3n} \beta_m x^m = \sum_{m=0}^{3n} \sum_{r=m}^{3n} \binom{r}{m} \alpha_r x^m \\ &= \sum_{r=0}^{3n} \alpha_r \sum_{m=0}^r \binom{r}{m} x^m = \sum_{r=0}^{3n} \alpha_r (x+1)^r = A(x+1). \end{aligned}$$

Equivalently, $A(x) = B(x-1)$, which is the principle of inclusion-exclusion, concisely stated. Expanding

$$B(x-1) = \sum_{m=0}^{3n} \beta_m (x-1)^m$$

and equating coefficients gives

$$\alpha_r = \sum_{m=r}^{3n} (-1)^{m-r} \binom{m}{r} \beta_m.$$

Using equation (2), we get

$$\alpha_r = \sum_{m=r}^{3n} (-1)^{m-r} \binom{m}{r} \sum_{\substack{s+t+u+v+w=n \\ t+2u+2v+3w=m}} \binom{n}{s, t, u, v, w} \frac{(4n - m)!}{(4!)^s(2!)^t(2!)^v}.$$

Since $m = t + 2u + 2v + 3w$, and

$$(-1)^{m-r} = (-1)^{t+2u+2v+3w-r} = (-1)^{t+w-r},$$

we can eliminate m from the formula for α_r and write it as a single sum.

Theorem 2.

$$\alpha_r = \sum_{\substack{s+t+u+v+w=n \\ t+2u+2v+3w \geq r}} (-1)^{t+w-r} \binom{t+2u+2v+3w}{r} \times \binom{n}{s, t, u, v, w} \frac{(4s+3t+2u+2v+w)!}{(4!)^s (2!)^t (2!)^v}.$$

Corollary 3. The number of permutations with no matches is

$$\alpha_0 = \sum_{s+t+u+v+w=n} (-1)^{t+w} \binom{n}{s, t, u, v, w} \frac{(4s+3t+2u+2v+w)!}{(4!)^s (2!)^t (2!)^v}.$$

Remark. The values of α_0 are given by the sequence A321633 in the OEIS. For $0 \leq n \leq 5$ the terms are

$$1, \quad 0, \quad 2, \quad 1092, \quad 2265024, \quad 11804626080.$$

Corollary 4. For a deck with four suits and n cards in each suit, we have

$$P(M = r) = \frac{\alpha_r}{(4n)! / (24)^n}, \quad r = 0, 1, \dots, 3n.$$

Example 5. For $k = 4$ and $n = 2$, there are 70 permutations of $\{1, 1, 1, 1, 2, 2, 2, 2\}$. Thus, $\beta_0 = 70$. It is easy to see that $\beta_6 = 2$. We will find β_3 and leave the rest as an exercise. To force at least three matches, the two ranks, denote them x and y , fit one of three patterns

$$\begin{aligned} xxxx \cdot y \cdot y \cdot y \cdot y & \quad s = 1, t = u = v = 0, w = 1, \\ xxx \cdot x \cdot yy \cdot y \cdot y & \quad s = 0, t = 1, u = 1, v = 0, w = 0, \\ xx \cdot xx \cdot yy \cdot y \cdot y & \quad s = 0, t = 1, u = 0, v = 1, w = 0. \end{aligned}$$

For each pattern, the multinomial coefficient $\binom{n}{s, t, u, v, w} = 2$, and so

$$\beta_3 = 2 \left(\frac{5!}{4!} + \frac{5!}{2!} + \frac{5!}{2!2!} \right) = 190.$$

After calculating the remaining β_m , you will get

$$B(x) = 70 + 210x + 270x^2 + 190x^3 + 78x^4 + 18x^5 + 2x^6.$$

Then

$$\begin{aligned} A(x) &= 70 + 210(x-1) + 270(x-1)^2 + 190(x-1)^3 \\ &\quad + 78(x-1)^4 + 18(x-1)^5 + 2(x-1)^6 \\ &= 2 + 6x + 18x^2 + 18x^3 + 18x^4 + 6x^5 + 2x^6. \end{aligned}$$

An easy partial check is to list the six permutations with exactly one match.

2. THE GENERAL CASE (k SUITS)

The patterns in which the k cards of a given rank can be grouped correspond to partitions of k . As we have seen, there are five partitions for $k = 4$. We represent a partition of k by a vector $\pi = (\pi_1, \pi_2, \dots, \pi_k)$, where π_i is the number of occurrences of i . Therefore, $\sum_i i\pi_i = k$ and $0 \leq \pi_i \leq k$. Let $\nu(\pi) = \pi_1 + \dots + \pi_k$ denote the number of parts in π , and let $\mathcal{P}(k)$ be the set of partitions of k .

Suppose that for each rank 1 to n , we choose a partition describing the grouping of the cards in that rank. Let s_π be the number of ranks which follow the partition π . Thus, $\sum_{\pi \in \mathcal{P}(k)} s_\pi = n$. Treating each group as a single object means that for a rank whose cards are grouped according to π , there are $\nu(\pi)$ objects to be permuted. The total number of objects to be permuted is $\sum_\pi s_\pi \nu(\pi)$, and so the factorial of this is the total number of permutations. But we need to divide by factors to account for identical symbols, i.e., for card groups of the same size in a rank. Thus, for each partition π we have a factor $\pi_i!$ repeated s_π times. Therefore, with this choice of a partition for each rank, we generate a list of permutations of size

$$\frac{\left(\sum_{\pi \in \mathcal{P}(k)} s_\pi \nu(\pi) \right)!}{\prod_{\pi \in \mathcal{P}(k)} \prod_{i=1}^k (\pi_i!)^{s_\pi}}.$$

It may help to examine the situation with $k = 4$ in order to see that this formula agrees with equation (1). There the numbers s, t, u, v, w are the values s_π for $\pi = (4, 0, 0, 0), (2, 1, 0, 0), (1, 0, 1, 0), (0, 2, 0, 0), (0, 0, 0, 1)$.

| s_π | π | π | $\nu(\pi)$ |
|---------|-----------------------------|----------------|------------|
| s | $x \cdot x \cdot x \cdot x$ | $(4, 0, 0, 0)$ | 4 |
| t | $xx \cdot x \cdot x$ | $(2, 1, 0, 0)$ | 3 |
| u | $xxx \cdot x$ | $(1, 0, 1, 0)$ | 2 |
| v | $xx \cdot xx$ | $(0, 2, 0, 0)$ | 2 |
| w | $xxxx$ | $(0, 0, 0, 1)$ | 1 |

The numerator is $(4s + 3t + 2u + 2v + w)!$. The denominator is the product of powers of the factorials of all the entries of all the π . Many of them are powers of $0!$ and $1!$, which we do not write, and so we get only three interesting factors: $(4!)^s (2!)^t (2!)^v$.

To define β_m , we construct permutations that have at least m matches. Grouping the cards in a single rank according to the partition π guarantees at least $\sum_i \pi_i(i - 1)$ matches because each part of size i contributes $i - 1$ matches. Let

$$\mu(\pi) = \sum_i \pi_i(i - 1).$$

The number of matches for a permutation with type (s_π) is at least $\sum_\pi s_\pi \mu(\pi)$. The number of ways to assign the patterns to suits to in order to have type (s_π) , where $\sum_\pi s_\pi = n$, is

the multinomial coefficient $\binom{n}{(s_\pi)}$. Define β_m by

$$\beta_m = \sum_{\substack{\sum_\pi s_\pi = n \\ \sum_\pi s_\pi \mu(\pi) = m}} \binom{n}{(s_\pi)} \frac{(\sum_{\pi \in \mathcal{P}(k)} s_\pi \nu(\pi))!}{\prod_{\pi \in \mathcal{P}(k)} \prod_{i=1}^k (\pi_i!)^{s_\pi}}.$$

Immediately from the definition of $\nu(\pi)$ and $\mu(\pi)$, we see that $\nu(\pi) + \mu(\pi) = n$, and so

$$\sum_{\pi \in \mathcal{P}(k)} s_\pi \nu(\pi) = \sum_{\pi \in \mathcal{P}(k)} s_\pi (n - \mu(\pi)) = \sum_{\pi \in \mathcal{P}(k)} s_\pi n - \sum_{\pi \in \mathcal{P}(k)} s_\pi \mu(\pi) = kn - m.$$

Therefore,

$$\beta_m = \sum_{\substack{\sum_\pi s_\pi = n \\ \sum_\pi s_\pi \mu(\pi) = m}} \binom{n}{(s_\pi)} \frac{(kn - m)!}{\prod_{\pi \in \mathcal{P}(k)} \prod_{i=1}^k (\pi_i!)^{s_\pi}}.$$

Let

$$B(x) = \sum_{m=0}^{(k-1)n} \beta_m x^m \quad \text{and} \quad A(x) = \sum_{r=0}^{(k-1)n} \alpha_r x^r,$$

where α_r is the number of permutations with exactly r matches. Then $A(x) = B(x - 1)$, and so

$$\alpha_r = \sum_{m=r}^{(k-1)n} (-1)^{m-r} \binom{m}{r} \beta_m.$$

3. BINOMIAL AND POISSON APPROXIMATION

The random variable M is the sum of identically distributed Bernoulli random variables, but they are not independent. The dependence, however, is rather weak: whether or not cards 1 and 2 match does not have much influence on whether or not cards 9 and 10 match. So, we consider a related random variable, call it M' , which is the sum of $kn - 1$ independent Bernoulli random variables with success probability $p = (k - 1)/(kn - 1)$. Thus, M and M' have the same expected value, and M' has a binomial distribution, which is much easier to compute:

$$P(M' = r) = \binom{kn - 1}{r} p^r (1 - p)^{kn - 1 - r}.$$

As we will see later (Corollary 11) it turns out that—somewhat surprisingly—they also have the same variance. Can we use M' as a good approximation to M ?

There is another approximation to consider. When n is large, the probability p of a match at any single site is small, but the expected number of matches is $k - 1$, independent of n . This suggests that a Poisson random variable with parameter $\lambda = k - 1$ may also be a good approximation. Recall that a Poisson random variable X with parameter $\lambda > 0$ has a distribution given by

$$P(X = k) = e^{-\lambda} \frac{\lambda^k}{k!} \quad \text{for } k = 0, 1, 2, \dots$$

Table 1 shows the computer calculations for the distributions of M , M' , and X for the standard deck where $n = 13$ and $k = 4$. For M , we use the formulas from Theorem 2 and Corollary 4. See the appendix for the code in Mathematica and Sage. The probabilities are rounded to the fifth decimal place. For $r \geq 14$ the probabilities round to 0.

TABLE 1. M for $k = 4, n = 13, M' \sim \text{Binomial}(51, 1/17), X \sim \text{Poisson}(3)$.

| r | $P(M = r)$ | $P(M' = r)$ | $P(X = r)$ |
|-----|------------|-------------|------------|
| 0 | 0.04548 | 0.04542 | 0.04979 |
| 1 | 0.14477 | 0.14477 | 0.14936 |
| 2 | 0.22611 | 0.22620 | 0.22404 |
| 3 | 0.23085 | 0.23091 | 0.22404 |
| 4 | 0.17321 | 0.17319 | 0.16803 |
| 5 | 0.10181 | 0.10175 | 0.10082 |
| 6 | 0.04879 | 0.04875 | 0.05041 |
| 7 | 0.01959 | 0.01959 | 0.0216 |
| 8 | 0.00672 | 0.00673 | 0.0081 |
| 9 | 0.00200 | 0.00201 | 0.0027 |
| 10 | 0.00052 | 0.00053 | 0.00081 |
| 11 | 0.00012 | 0.00012 | 0.00022 |
| 12 | 0.00002 | 0.00003 | 0.00006 |
| 13 | 0.00000 | 0.00000 | 0.00001 |

For all practical purposes, the distributions of M and M' are the same. To measure how far apart they are, we use the *total variation distance*

$$\begin{aligned} d_{TV}(M, M') &:= \frac{1}{2} \sum_{r=0}^{51} |P(M = r) - P(M' = r)| \\ &= 0.000181682. \end{aligned}$$

An equivalent formulation of the total variation distance is

$$d_{TV}(M, M') = \sup_E |P(M \in E) - P(M' \in E)|,$$

where E ranges over all events, which in this case are the subsets of $\{0, 1, 2, \dots, 51\}$. Thus, the maximal difference between the probabilities on any event is less than two hundredths of one percent.

As n increases, the dependence among the M_i decreases, and so we expect that M and M' become ever closer. The limit distribution of M' is Poisson with $\lambda = 3$, and so the same should hold for M . We can prove that by showing that the total variation distance between the distributions of M and M' goes to 0. To do that, we estimate the total variation distance

using a result of Soon [3, Corollary 1.5], which involves all the covariances $\text{Cov}(M_i, M_j)$. Soon's estimate, applied to our situation, says that

$$(3) \quad d_{TV}(M, M') \leq C_{k,n} \sum_{j \neq i} |\text{Cov}(M_i, M_j)|,$$

where

$$C_{k,n} = \frac{1 - p^{kn} - (1 - p)^{kn}}{kn p (1 - p)} \quad \text{and} \quad p = \frac{k - 1}{kn - 1}.$$

It is straightforward to show that $\lim_{n \rightarrow \infty} C_{n,k} = 1/(k - 1)$. The covariance of M_i and M_j is given by

$$\text{Cov}(M_i, M_j) = E(M_i M_j) - E(M_i)E(M_j).$$

We have $E(M_i)E(M_j) = (k - 1)^2/(kn - 1)^2$. The calculation of $\text{Cov}(M_i M_j)$ is carried out through the following three lemmas.

Lemma 6. If $|i - j| = 1$, then

$$E(M_i M_j) = \frac{(k - 1)(k - 2)}{(kn - 1)(kn - 2)}.$$

Proof. We may assume $j = i + 1$. Then $E(M_i M_j)$ is the probability that cards in location i , $i + 1$, and $i + 2$ are the same rank. The first card can be anything. The probability that the next card has the same rank is $(k - 1)/(kn - 1)$, and then the probability that the third card has the same rank is $(k - 2)/(kn - 2)$. \square

Lemma 7. If $|i - j| > 1$, then

$$E(M_i M_j) = \frac{(k - 1)(k - 2)(k - 3) + (k - 1)^2(kn - k)}{(kn - 1)(kn - 2)(kn - 3)}.$$

Proof. In this case, the locations of the matches are separated. There are two possible configurations, namely

$$\dots xx \dots xx \dots \quad \text{or} \quad \dots xx \dots yy \dots$$

In the first configuration, there are four cards of the same rank; in the second the matches involve cards of different ranks. The probability of the first is

$$\frac{(k - 1)(k - 2)(k - 3)}{(kn - 1)(kn - 2)(kn - 3)}.$$

Note that this configuration cannot occur for $k \leq 3$, and the formula correctly gives 0 in that case. For the second configuration, the probability is

$$\frac{(k - 1)^2(kn - k)}{(kn - 1)(kn - 2)(kn - 3)}$$

because $(k - 1)/(kn - 1)$ is the probability of the xx match. Then y must be one of $kn - k$ cards of the remaining $kn - 2$ and the second y must be one of $k - 1$ cards of the remaining $kn - 3$. The sum of the probabilities for the two configurations is $E(M_i M_j)$. \square

Lemma 8.

$$\text{Cov}(M_i, M_j) = \begin{cases} \frac{(k-1)(k-kn)}{(kn-1)^2(kn-2)} & \text{if } |i-j| = 1, \\ \frac{2(k-1)(kn-k)}{(kn-1)^2(kn-2)(kn-3)} & \text{if } |i-j| > 1. \end{cases}$$

Proof. Routine algebra left to the reader. \square

Notice that the covariance is negative when $|i-j| = 1$ and positive when $|i-j| > 1$.

Theorem 9. *The total variation distance between the distributions of M and M' goes to 0 as $n \rightarrow \infty$.*

Proof. In the estimate given in equation (3), the constant $C_{k,n}$ has a limit and so it is bounded. We need to show that $\sum_{i \neq j} |\text{Cov}(M_i, M_j)| \rightarrow 0$ as $n \rightarrow \infty$. Break the sum into two pieces. There are $2(kn-2)$ terms with $|i-j| = 1$ and the covariances are negative. From the previous lemma we get

$$\begin{aligned} \sum_{|i-j|=1} |\text{Cov}(M_i, M_j)| &= -2(kn-2) \frac{(k-1)(k-kn)}{(kn-1)^2(kn-2)} \\ &= \frac{2(k-1)(kn-k)}{(kn-1)^2}. \end{aligned}$$

There are $(kn-2)(kn-3)$ terms with $|i-j| > 1$, and the covariances are positive. Again from the lemma we get

$$\begin{aligned} \sum_{|i-j|>1} |\text{Cov}(M_i, M_j)| &= (kn-2)(kn-3) \frac{(k-1)(2kn-2k)}{(kn-1)^2(kn-2)(kn-3)} \\ &= \frac{2(k-1)(kn-k)}{(kn-1)^2}. \end{aligned}$$

It follows that

$$\lim_{n \rightarrow \infty} \sum_{i \neq j} |\text{Cov}(M_i, M_j)| = \lim_{n \rightarrow \infty} \frac{4(k-1)(kn-k)}{(kn-1)^2} = 0.$$

\square

Corollary 10. *As $n \rightarrow \infty$ the distribution of M converges to a Poisson distribution with parameter $\lambda = k-1$. That is, for $r \geq 0$,*

$$\lim_{n \rightarrow \infty} \text{P}(M = r) = e^{k-1} \frac{(k-1)^r}{r!}.$$

Proof. From the theorem, we know that for $r \geq 0$, we have

$$\lim_{n \rightarrow \infty} |\text{P}(M = r) - \text{P}(M' = r)| = 0.$$

Thus, $\text{P}(M = r)$ and $\text{P}(M' = r)$ have the same limit if the limit exists. In fact, the limit does exist because the distribution of M' converges to the Poisson distribution with parameter

$k - 1$. This is a classical result which says that for a sequence $p_N > 0$ such that $Np_N \rightarrow \lambda$, the binomial distribution with parameters N and p_N converges to the Poisson distribution with parameter λ . This means that for each $r \geq 0$,

$$\lim_{N \rightarrow \infty} \binom{N}{r} p_N^r (1 - p_N)^{N-r} = e^{-\lambda} \frac{\lambda^r}{r!}.$$

For M' we have $N = kn - 1$, $p_N = (k - 1)/(kn - 1)$, and $\lambda = k - 1$. \square

We have noted the curious fact that M and M' have the same variance. A more intuitive explanation would be of interest, but our proof is simply the computation that depends on the fact that $\sum_{|i-j|=1} \text{Cov}(M_i, M_j) = -\sum_{|i-j|>1} \text{Cov}(M_i, M_j)$.

Corollary 11.

$$\text{Var}(M) = \text{Var}(M') = \frac{k(k-1)(n-1)}{kn-1}.$$

Proof.

$$\begin{aligned} \text{Var}(M) &= \text{Var}\left(\sum_i M_i\right) = \sum_i \text{Var}(M_i) + \sum_{i \neq j} \text{Cov}(M_i, M_j) \\ &= \sum_i \text{Var}(M_i) + \sum_{|i-j|=1} \text{Cov}(M_i, M_j) + \sum_{|i-j|>1} \text{Cov}(M_i, M_j) = \sum_i \text{Var}(M_i). \end{aligned}$$

Since M' is the sum of *independent* Bernoulli random variables with the same distribution as the M_i , we also have $\text{Var}(M') = \sum_i \text{Var}(M_i)$. The M_i are Bernoulli random variables with parameter $p = (k - 1)/(kn - 1)$ and variance $p(1 - p)$. Thus,

$$\begin{aligned} \sum_i \text{Var}(M_i) &= \sum_i p(1 - p) = (kn - n)p(1 - p) \\ &= \frac{k(k-1)(n-1)}{kn-1}. \end{aligned}$$

\square

4. TWO RANKS ($n = 2$)

For n even moderately large, the binomial approximation is quite good, but it is not so good for small values of n , especially for $n = 2$. In that case, however, counting permutations according to the number of matches is quite tractable for general k , and it does not use the general formula developed in section 5, which becomes unwieldy as k increases.

For $n = 2$ and $k \geq 1$, the number of permutations is $\binom{2k}{k}$. If a permutation begins with 1, then it is made up of a string of 1's, followed by a string 2's, then a string of 1's, and so on.

Suppose s is even, say $s = 2t$. Then the permutation is made up of t strings of 1's and t strings of 2's that are interlaced, and so it is completely determined by two increasing sequences

$$1 \leq u_1 < u_2 < \cdots < u_{t-1} \leq k - 1, \quad 1 \leq v_1 < v_2 < \cdots < v_{t-1} \leq k - 1$$

where u_i is the total number of 1's that are in the first i strings of 1's and v_i is the total number of 2's in the first i strings of 2's. For example, if $k = 6$ and the permutation is 112111222122, then $t = 3$, $u_1 = 2, u_2 = 5$ and $v_1 = 1, v_2 = 4$. The number of permutations is $\binom{k-1}{t-1}^2$.

Suppose s is odd, say $s = 2t + 1$. Then there are $t + 1$ strings of 1's and t strings of 2's so that the permutation is determined by sequences

$$1 \leq u_1 < u_2 < \cdots < u_t \leq k - 1, \quad 1 \leq v_1 < v_2 < \cdots < v_{t-1} \leq k - 1.$$

Therefore, the number of these permutations is $\binom{k-1}{t} \binom{k-1}{t-1}$.

Theorem 12. *If $n = 2$, $k \geq 1$, and $0 \leq r \leq 2k - 2$, then α_r , the number of permutations with exactly r matches, is given by*

$$\alpha_r = \begin{cases} 2 \binom{k-1}{k-\ell-1}^2 & \text{if } r = 2\ell, \\ 2 \binom{k-1}{k-\ell-1} \binom{k-1}{k-\ell-2} & \text{if } r = 2\ell + 1. \end{cases}$$

Proof. Above we have counted the permutations beginning with 1 and made up of s strings. In a deck of $2k$ cards there are $2k - 1$ locations where there could be a match, but each boundary between two strings reduces the number of matches by one. There are $s - 1$ boundaries between strings. Therefore, $r = 2k - 1 - (s - 1) = 2k - s$. If r is even, say $r = 2\ell$, then s is even since $r = 2k - s$. Letting $s = 2t$ we see that $t = k - \ell - 1$. Thus, the number of permutations beginning with 1 and having r matches is

$$\binom{k-1}{t-1}^2 = \binom{k-1}{k-\ell-1}^2.$$

If r is odd, say $r = 2\ell + 1$, then s is odd, say $s = 2t + 1$, and $t = k - \ell - 1$. The number of permutations beginning with 1 and having r matches is

$$\binom{k-1}{t} \binom{k-1}{t-1} = \binom{k-1}{k-\ell-1} \binom{k-1}{k-\ell-2}.$$

Now we double these numbers because there are just as many permutations that start with 2. □

You may have noticed in Example 5 that the generating function $A(x)$ is palindromic. That is, $\alpha_r = \alpha_{2k-2-r}$. That can be proved in general for any k as long as $n = 2$ either from the formula in Proposition 12 or from the correspondence between permutations and the sequences (u_i) and (v_i) . The second approach replaces each sequence by its complement and has the virtue of giving a bijection between permutations with r matches and those with $2k-2-r$ matches. The bijection is not at all obvious. In the case that $r = 2k-2-r$, the bijection is not the identity map. The palindromic property does not continue for $n \geq 3$. In

particular, although $\alpha_{kn-n} = n!$, you can see that α_0 , the number of permutations with no matches, is larger simply by constructing enough of them.

The numbers α_r for $k \geq 1$ and $0 \leq r \leq 2k - 2$ appear as entry A152659 in the OEIS, but there they are described as the number of lattice paths from $(0, 0)$ to (n, n) with steps $(1, 0)$ and $(0, 1)$ and having k turns. To see the correspondence between permutations and paths, you must replace n with k and replace k with $2k - r - 1$.

APPENDIX: MATHEMATICA AND SAGE

The Mathematica code below samples 100,000 random deals and counts the number of matches in each deal. Then it finds the average and plots a histogram. The deck is the multiset with four copies of each number from 1 to 13.

```
deck = Flatten[Table[{i, i, i, i}, {i, 1, 13}]];
matches = {};
Do[s = RandomSample[deck, 52];
  AppendTo[matches, Count[Differences[s], 0]], {100000}];
Mean[matches]
Histogram[matches, {1}, "Probability",
  Ticks->{Range[0, 13], Automatic}]
```

Mathematica code for β_m and α_r with n ranks and four suits.

```
Needs["Combinatorica`"]
beta[n_, m_] :=
Module[{z=Select[Compositions[n, 5],
  #[[2]]+2#[[3]]+2#[[4]]+3#[[5]]==m&}},
  Sum[Multinomial[z[[i, 1]], z[[i, 2]], z[[i, 3]], z[[i, 4]], z[[i, 5]]]
  ((4, 3, 2, 2, 1).z[[i]])!/(24^z[[i, 1]]*2^(z[[i, 2]]+z[[i, 4]])),
  {i, 1, Length[z]}]]
alpha[n_, r_] := Sum[(-1)^m beta[n, m], {m, r, 3n}]
```

Sage code for the random samples and histogram.

```
def differences(x):
    return [x[i+1]-x[i] for i in range(len(x)-1)]
deck = 4*[1..13]
matches = []
for i in range(100000):
    x = Permutations(deck).random_element()
    matches.append(differences(x).count(0))
import numpy
```

```
import matplotlib.pyplot as plt
plt.figure(figsize=(7.5, 5))
plt.hist(matches, [i-.5 for i in range(13)], density=True, \
           edgecolor="white")
plt.xticks(range(14))
plt.show()
```

Sage code for β_m and α_r with n ranks and four suits.

```
def beta(n, m):
    cmp=Compositions(n, length=5, min_part=0)
    z=[c for c in cmp if c[1]+2*c[2]+2*c[3]+3*c[4]==m]
    return sum(multinomial(list(x))
               *factorial(4*x[0]+3*x[1]+2*x[2]+2*x[3]+x[4])
               /((24)^x[0]*2^x[1]*2^x[3]) for x in z)
def alpha(n, r):
    return sum((-1)^(m-r)*binomial(m, r)*beta(n, m) \
               for m in range(r, 3*n+1))
```

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